

Similar to Hilbert algebras, commutative L-algebras can be derived from logic. Both are BCK-algebras. Their logic is given by the axioms (25)

$$(B) \quad (p \cdot q) \cdot ((r \cdot p) \cdot (r \cdot q))$$

$$(C) \quad (p \cdot (q \cdot r)) \cdot (q \cdot (p \cdot r))$$

$$(K) \quad p \cdot (q \cdot p)$$

and the rule (MP)  $p, p \cdot q \vdash q$ .

Using (C), axiom (B) is equivalent to

$$(B') \quad (r \cdot p) \cdot ((p \cdot q) \cdot (r \cdot q)).$$

Hence

$$p \leq q : \iff \vdash p \cdot q$$

is transitive:  $r \leq p \leq q \xrightarrow{(B')} \vdash (p \cdot q) \cdot (r \cdot q) \vdash r \cdot q \Rightarrow r \leq q$ . Furthermore,

(C) with  $r=p$  gives  $\vdash (p \cdot (q \cdot p)) \cdot (q \cdot (p \cdot p)) \xrightarrow{(K)} q \cdot (p \cdot p) \vdash p \cdot p$ . So  $p \leq p$ , and

$$p \equiv q : \iff p \leq q \leq p$$

is an equivalence relation. By (B) and (B'), it is a congruence relation.

So, again, we can form the Lindenbaum algebra over the variable set  $V$ :

$$L(V) := P / \equiv$$

with the greatest element 1, the class of true formulae. Now  $L(V)$  satisfies

$$(B) \quad p \cdot q \leq (r \cdot p) \cdot (r \cdot q)$$

$$(C) \quad p \cdot (q \cdot r) = q \cdot (p \cdot r)$$

$$(K) \quad p \leq q \cdot p$$

and  $1 = p \cdot p = q \cdot 1$  (since  $\vdash q \cdot (p \cdot p)$ ). By (C), we have

$$p \leq q \cdot r \iff q \leq p \cdot r.$$

Hence  $1 \cdot p \leq 1 \cdot p$  gives  $1 \in (1 \cdot p) \cdot p$ . Thus  $1 \cdot p \leq p \leq 1 \cdot p$ . So  $L(V)$  is a BCK-algebra,

i.e. BCK-algebras are the remainders of BCK-logic.

Prop. 31. An L-algebra  $X$  satisfying (C) is a BCK-algebra.

Pr. For  $r=p$ , (C) gives  $p \cdot (q \cdot p) = q \cdot (p \cdot p) = q \cdot 1 = 1$ .  $\xrightarrow{(K)} p \cdot q \leq$

$(q \cdot r) \cdot (p \cdot q) = (r \cdot p) \cdot (r \cdot q)$ .  $\rightarrow$  (B).  $\square$

So we could speak of a BCKL-algebra or a CL-algebra. We choose a middle course and call these algebras CKL-algebras. They are also known as HBCK-algebras.

It was conjectured by Woroniski (1985) that CKL-algebras form a variety (proved by Blok and Ferreirim 2001). Indeed, the axiom  $x \cdot y = y \cdot x = 1 \Rightarrow x = y$  can be replaced by Cornish's identity

$$(J) \quad (((x \cdot y) \cdot y) \cdot x) \cdot x = (((y \cdot x) \cdot x) \cdot y) \cdot y.$$

Prop. 32. Let  $X$  be an  $L$ -algebra.  $\square$

- (a)  $X$  is a CKL-algebra
- (b)  $x \leq y \cdot z \Rightarrow y \leq x \cdot z$
- (c)  $X$  is a KL-algebra with  $x \leq (x \cdot y) \cdot y$  for all  $x, y \in X$ .

Pr. (a)  $\Rightarrow$  (b) is trivial.

(b)  $\Rightarrow$  (c):  $x \cdot y \leq x \cdot y \cdot z \rightarrow x \leq (x \cdot y) \cdot y$ . For  $z = y$ , we get (K).

(c)  $\Rightarrow$  (a):  $a := (x \cdot (y \cdot z)) \cdot (x \cdot z) = ((y \cdot z) \cdot x) \cdot ((y \cdot z) \cdot z) \stackrel{(K)}{\geq} (y \cdot z) \cdot z \geq y$ .  $\rightarrow$   
 $x \cdot (y \cdot z) \leq a \cdot (x \cdot z) \stackrel{(K)}{\leq} (a \cdot y) \cdot (a \cdot (x \cdot z)) = (y \cdot a) \cdot (y \cdot (x \cdot z)) = y \cdot (x \cdot z), \square$

Cor. Every commutative  $L$ -algebra  $X$  is a CKL-algebra.

Pr. By Prop. 8,  $X$  satisfies (K) and (V):  $(x \cdot y) \cdot y = (y \cdot x) \cdot x (= x \vee y)$ . So  $x \stackrel{(K)}{\leq} (y \cdot x) \cdot x = (x \cdot y) \cdot y. \square$

Recall that an MV-algebra is a commutative  $L$ -algebra with  $0$ . With

$$x' := x \cdot 0$$

every MV-algebra has a multiplication

$$x \cdot y := (x \cdot y')', \text{ i.e. } x \cdot y = (x \cdot y')', \text{ (in } \mathcal{L}(X): (x \cdot y')' = x \vee y)$$

Prop. 33. Every MV-algebra  $X$  is a commutative monoid and a lattice.

with  $x'' = x$  and

$$\begin{aligned} x' \cdot y' &= y \cdot x \\ x \cdot y \cdot z &= x \cdot (y \cdot z) \\ x \wedge y &= (x \cdot y) \cdot x; \quad x \vee y = (x \cdot y) \cdot y \end{aligned}$$

Pr.  $x'' = (x \cdot 0) \cdot 0 = (0 \cdot x) \cdot x = 1 \cdot x = x. \rightarrow x' \cdot y' = x' \cdot (y \cdot 0) = y \cdot (x' \cdot 0) = y \cdot x. \rightarrow$   
 $x \cdot y = (x \cdot y')' = (x'' \cdot y'')' = (y \cdot x')' = y \cdot x,$  and  $x \cdot y \cdot z = (x \cdot y')' \cdot z = z' \cdot (x \cdot y') = x \cdot (z' \cdot y') = x \cdot (y \cdot z). \rightarrow (x \cdot y) \cdot z = (x \cdot y \cdot z)' = (x \cdot (y \cdot z'))' = (x \cdot (y \cdot z))' = x \cdot (y \cdot z).$  Furthermore,  
 $1 \cdot x = (1 \cdot x')' = x'' = x.$  So  $X$  is a monoid. The equation  $x' \cdot y' = y \cdot x$  implies that  
 $x' \leq y' \Leftrightarrow y \leq x.$  So  $x \wedge y = (x' \vee y')' = ((x' \cdot y') \cdot y')' = (x' \cdot y') \cdot y = (y \cdot x) \cdot y, \square$

Next we show that the elements of an MV-algebra can be regarded as measurable sets. To say that two sets  $A, B$  have the same measure means to split  $A$  into small pieces, translate them with a group and reassemble them to  $B$ .

Def. A measure  $\mu: X \rightarrow Y$  between  $L$ -algebras  $X, Y$  is a map with  
 $x \geq y \implies \mu(x \cdot y) = \mu(x) \cdot \mu(y), \mu(x) \geq \mu(y).$

Thus  $x=y$  gives  $\mu(1) = 1$ . With  $X \subset S(X)$  we get

Prop. 34.  $\mu: X \rightarrow Y$  is a measure iff  
 $\exists xy \in X \implies \mu(xy) = \mu(x)\mu(y).$

Pf. In  $S(X)$ :  $xy \leq z \iff x \leq y \cdot z$ , and  $z \leq xy \iff (y \cdot z \cdot x)(z \cdot y) = 1 \iff y \cdot z \leq x$   
and  $z \leq y$ . So

$$xy = z \iff (x = y \cdot z \text{ and } y \geq z).$$

So the condition of Prop. 34 is equivalent to

$$x = y \cdot z, y \geq z \implies \mu(x) = \mu(y) \cdot \mu(z), \mu(y) \geq \mu(z). \quad \square$$

In additive terms:  $\exists x+y \implies \mu(x+y) = \mu(x) + \mu(y)$ . (finite measure)

Now let  $X$  be an MV-algebra. By Thm. 3,  $G(X)$  is a commutative  $l$ -group with  $X = [0, 1] \subset S(X) = G(X)^+$ , and  $x \vee y = (x \cdot y) \cdot y$ .

Cor. Let  $X, Y$  be MV-algebras. Every measure  $\mu: X \rightarrow Y$  extends uniquely to a group homomorphism  $G(\mu): G(X) \rightarrow G(Y)$ , and every group homomorphism  $f: G(X) \rightarrow G(Y)$  with  $f(X) \subset Y$  restricts to a measure  $\mu: X \rightarrow Y$ .

Pf. Let  $\mu: X \rightarrow Y$  be a measure. Then

$$(1) \quad \mu(ab) = \mu(a) \mu(b)$$

holds for  $0 \leq ab$  in  $S(X)$ . Every  $a \in S(X)$  satisfies  $0^n \leq a$  for some  $n \in \mathbb{N}$ . If  $0^{n+1} \leq a$ , then  $0^n \leq 0 \cdot a$ , and  $a \cdot 0 = (a \cdot 0)'(a \cdot 0) \cdot a = (a \cdot 0)'(0 \cdot a) \cdot 0$ .

$$a = (a \cdot 0)'(0 \cdot a) \text{ with } (a \cdot 0)' \in X.$$

Thus, recursively, we define

$$(2) \quad \mu(a) := \mu((a \cdot 0)') \mu(0 \cdot a).$$

Assume that (1) holds for  $0^n \leq a, b \leq 1$ . We show that for  $x, y \in X$  with  $xa = yb$ ,

$$(3) \quad \mu(x) \mu(a) = \mu(y) \mu(b).$$

Now  $xa = yb \iff z := (a \cdot b) \cdot x = (b \cdot 0) \cdot y$  with  $x \leq a \cdot b, y \leq b \cdot a$ . So  $z(a \cdot b) = a \cdot b \wedge x = x$ , and similarly,  $z(b \cdot a) = y$ .  $\implies \mu(z) \mu(a \cdot b) = \mu(x), \mu(z) \mu(b \cdot a) = \mu(y)$ . So we have to prove:  $\mu(a \cdot b) \mu(a) \stackrel{!}{=} \mu(b \cdot a) \mu(b)$ , which holds by induction.

Now we extend (3) to  $0^{n+1} \leq xa$ , that is,  $\mu(x) \mu(a) \stackrel{!}{=} \mu(xa)$ . By (2), i.e.

$$\mu(x) \mu((a \cdot 0)') \mu(0 \cdot a) \stackrel{!}{=} \mu((xa \cdot 0)') \mu(0 \cdot xa)$$

Since  $y := a \cdot 0 \in X$ ,  $0 \cdot x a = (y \cdot x)(0 \cdot a) \rightsquigarrow \mu(0 \cdot x a) = \mu(y \cdot x) \mu(0 \cdot a)$ . So

$$\mu(x) / \mu(y) = \mu(x \cdot y) / \mu(y \cdot x)$$

Indeed,  $x y' y = x 0 = 0 x = (x \cdot y)' (x \cdot y) x = (x \cdot y)' (y \cdot x) y \rightsquigarrow x y' = (x \cdot y)' (y \cdot x)$ .

The converse follows by Prop. 34.  $\square$

For an MV-algebra  $X$ , and  $x, y \in X \subset S(X)$ ,

$$xy \in X \iff 0 \leq xy \iff 0 \cdot xy = 1 \iff 1 = (y' \cdot x)(0 \cdot y) = y' \cdot x \rightsquigarrow$$

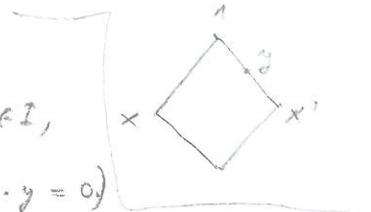
$$xy \in X \iff y' \leq x \iff x' \leq y$$

Example. If  $X$  is a Boolean algebra, then  $xy \in X \iff x' \leq y \iff x \vee y = 1$ .

Let  $I = [0, 1]$ , additively, with opposite ordering.  $\rightsquigarrow$

$S(I) = \mathbb{R}_{\geq 0}$ ,  $G(I) = \mathbb{R}$ , with logical unit 0, for  $x, y \in I$ ,

$x' = 1 - x$ , and  $x \cdot y = y - x$  if  $x \leq y$ . (For  $x \geq y$ , we have  $x \cdot y = 0$ )



Let  $\mathcal{B}(I)$  be the Boolean algebra generated by  $\mathcal{O}(I)$  (Borel sets).

So  $x + y \in I$  iff  $x + y \leq 1$ , and in  $\mathcal{B}(I)$ :  $A + B \in \mathcal{B}(I) \iff X \setminus A = B \iff A \cap B = \emptyset$ .

$\xrightarrow{\text{Prop. 34}}$  A measure  $\mu: \mathcal{B}(I) \rightarrow I$  is a map with

$$\mu(A \sqcup B) = \mu(A) + \mu(B)$$

Note that there are two extreme cases of MV-algebras: Boolean algebras (like  $\mathcal{B}(I)$ ) and MV-chains, i.e. linearly ordered MV-algebras (like  $I$ ).

Prop. 35. Let  $X$  be an MV-algebra, and  $I \subset X$ .  $\odot$

- (a)  $I$  is an (L-algebra) ideal.
- (b)  $1 \in I$ , and  $x, x \cdot y \in I \implies y \in I$ .
- (c)  $I$  is an ~~upper~~ upper set, and  $x, y \in I \implies xy \in I$
- (d)  $I \neq \emptyset$  is an upper set, and  $x, y \in I, x' \leq y \implies xy \in I$ .

If (a)  $\implies$  (b) is trivial, (b)  $\implies$  (a) follows since  $X$  is a CKL-algebra (Prop. 32).

(b)  $\implies$  (c): If  $x, y \in I$ , then  $y \cdot xy = y \cdot (x \cdot y)' = (x \cdot y)' \cdot y = x \vee y' \geq x \rightsquigarrow y, y', xy \in I$ .  $\rightsquigarrow xy \in I$ . (c)  $\implies$  (d) is trivial.

(d)  $\implies$  (b): If  $x, x \cdot y \in I$ , then  $y \geq x \wedge y = (x \cdot y) x \in I$ .  $\square$

For a measure  $\mu: X \rightarrow Y$ , we call

$$\text{Ker } \mu := \{x \in X \mid \mu(x) = 1\}$$

the kernel of  $\mu$ .

Prop. 36. Let  $\mu$  be a measure of MV-algebras. Then  $\text{Ker } \mu$  is an ideal, and  $\mu$  factors through  $X \rightarrow X / \text{Ker } \mu$ .



Prop 38. Let  $X$  be an MV-algebra.

- (a)  $y \mapsto x \cdot y$  is a lattice homomorphism  $X \rightarrow X$ .
- (b)  $x \mapsto x \cdot y$  is a lattice homomorphism  $X^{op} \rightarrow X$ .
- (c)  $X$  is a distributive lattice.

pf. We embed  $X$  into  $G(X)$ . Then  $x \cdot y = yx^{-1} \wedge 1$ .

a)  $x \cdot (y \vee z) = (y \vee z)x^{-1} \wedge 1 = yx^{-1} \wedge zx^{-1} \wedge 1 = x \cdot y \vee x \cdot z$ .

$$(x \vee y) \wedge (x \vee z) = (x \cdot y) \cdot y \wedge (x \cdot z) \cdot z \stackrel{(K)}{\leq} \left[ ((x \cdot y) \cdot (x \cdot z)) \cdot ((x \cdot y) \cdot y) \right] \wedge \left[ ((x \cdot z) \cdot (x \cdot y)) \cdot ((x \cdot z) \cdot z) \right]$$

$$\leq \underbrace{((x \cdot y) \cdot (x \cdot z)) \cdot ((x \cdot y) \cdot z)}_{(x \cdot y) \cdot (x \cdot z) \cdot (x \cdot y) \cdot z} \wedge \underbrace{((x \cdot z) \cdot (x \cdot y)) \cdot ((x \cdot z) \cdot z)}_{(x \cdot z) \cdot (x \cdot y) \cdot (x \cdot z) \cdot z}$$

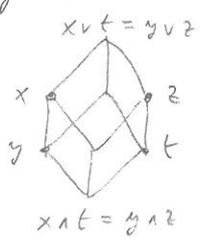
$$= (x \cdot y \wedge x \cdot z) \cdot y \wedge (x \cdot y \wedge x \cdot z) \cdot z = (x \cdot (y \vee z)) \cdot (y \vee z) = x \vee (y \vee z) \leq (x \vee y) \wedge (x \vee z) \text{ by (c)}$$

$$\rightarrow (x \cdot (y \vee z))x = x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z) = (x \cdot y)x \vee (x \cdot z)x = (x \cdot y \vee x \cdot z)x \text{ by (a)}$$

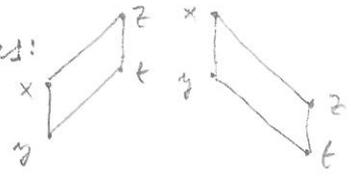
b)  $(x \vee y) \cdot z = z' \cdot (x \vee y)' = z' \cdot (x' \wedge y') = z' \cdot x' \wedge z' \cdot y' = x \cdot z \wedge y \cdot z$

$(x \wedge y) \cdot z = z' \cdot (x \wedge y)' = z' \cdot (x' \vee y') = z' \cdot x' \vee z' \cdot y' = x \cdot z \vee y \cdot z$   $\square$

Def. Let  $D$  be a distributive lattice. We call two intervals  $[y, x]$  and  $[t, z]$  equivalent if  $x \vee t = y \vee z$  and  $x \wedge t = y \wedge z$ :



special cases:

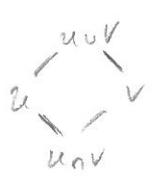


The equivalence class of  $[y, x]$  will be denoted by  $x/y$ .

Note: If  $X$  is an MV-algebra,  $x/y = z/t \iff x \cdot y = z \cdot t$ .

indeed,  $(x \vee y) \cdot y = ((x \cdot y) \cdot y) \cdot y = x \cdot y = x \cdot (x \wedge y)$ , where  $((x \cdot y) \cdot y) \cdot y = x \cdot y$  follows from  $x \leq (x \cdot y) \cdot y \rightarrow ((x \cdot y) \cdot y) \cdot y \leq x \cdot y \leq ((x \cdot y) \cdot y) \cdot y$ .

Let  $B(D)$  be the Boolean algebra generated by  $D$ . It can be realized as  $\mathcal{L}(X)$  if  $X = \text{Spec } D$  and thus  $D = \mathcal{S}(X)$ . For  $u, v \in \mathcal{S}(X)$  we have



$(u \vee v) \setminus v = u \setminus v = u \setminus (u \wedge v)$ . Conversely, assume that  $u_1 \supseteq u_2$  and  $v_1 \supseteq v_2$ . Then  $u_1 \setminus u_2 = v_1 \setminus v_2$  implies  $u_1 = u_2 \vee (u_1 \setminus u_2) = u_2 \vee (v_1 \setminus v_2)$ .  $\rightarrow u_1 \vee v_2 = u_2 \vee v_2$  and  $u_1 \wedge v_2 = u_2 \wedge v_2$ . Thus, by symmetry,

$u_1 \wedge v_2 = u_2 \wedge v_1$ . Hence

$$u_1/u_2 = v_1/v_2 \iff u_1 \setminus u_2 = v_1 \setminus v_2$$

Denote the complement of  $a \in B(D)$  by  $a^c$ . So we obtain for  $x_1 \geq x_2, y_1 \geq y_2$  in  $D$ :

$$x_1/x_2 = y_1/y_2 \iff x_1 \wedge x_2^c = y_1 \wedge y_2^c \text{ in } B(D).$$

Prop. 39. Let  $D$  be a distributive lattice. Every  $a \in B(D)$  is a disjoint union (39)

$$a = (x_1 \wedge y_1^c) \vee \dots \vee (x_n \wedge y_n^c), \quad x_i, y_i \in D. \quad (1)$$

Pr.  $(x \wedge y^c) \wedge (z \wedge t^c)^c = x \wedge y^c \wedge (z^c \vee t) = (x \wedge y^c \wedge z^c) \vee (x \wedge y^c \wedge t) = (x \wedge (y \vee z)^c) \vee (x \wedge t \wedge y^c)$ , disjoint since  $z^c \wedge t \leq z^c \wedge z = 0$ . Furthermore,

$$(x \wedge y^c) \wedge (z \wedge t^c) = (x \wedge z) \wedge (y \vee t)^c. \quad (2)$$

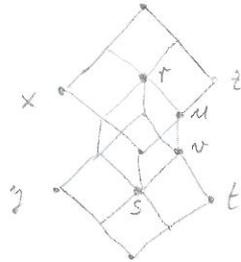
Thus, if  $a = a_1 \vee \dots \vee a_n$  is disjoint with  $a_i = x_i \wedge y_i^c$  and  $b = x \wedge y^c$ , then  $a \wedge b^c = (a_1 \wedge b^c) \vee \dots \vee (a_n \wedge b^c)$  is in normal form, and also  $a \vee b = (a \wedge b^c) \vee b$ .  $\square$

Replacing  $y_i$  by  $x_i \wedge y_i$ , we can assume that  $y_i < x_i$  (removing zero terms).

So the  $x_i \wedge y_i := x_i \wedge y_i^c$  correspond to quotients  $x_i/y_i$ . Eq. (2) gives

$$(x \wedge y) \wedge (z \wedge t) = 0 \iff x \wedge z \leq y \vee t.$$

The sublattice generated by  $x, y, z, t$  is



So  $u/v$  corresponds to  $(x \wedge y) \wedge (z \wedge t)$ .

We call  $x/y$  and  $z/t$  disjoint if

$x \wedge z \leq y \vee t$ . (i.e.  $u/v$  trivial).

$$u = (x \wedge z) \vee t; \quad v = (y \wedge z) \vee t.$$

If  $x/y$  and  $z/t$  are disjoint, we write  $(x/y) + (z/t) := (x \wedge y) \vee (z \wedge t)$ .

Accordingly, we have a min  $x/y + z/t$  of disjoint quotients. So (1) can be written as

$$a = (x_1 \wedge y_1) + \dots + (x_n \wedge y_n). \quad (3)$$

If  $x/y$  and  $z/t$  are disjoint, we have  $x/y + z/t = s/(y \wedge t) + (x \vee z)/t$ , where  $y \wedge t \leq s \leq r \leq x \vee z$ .

By induction, every min (3) can be put into a form where  $x_1 > y_1 \geq x_2 > y_2 \geq \dots$ . So

Cor. Let  $D$  be a distributive lattice. Every  $a \in B(D)$  is a disjoint union

$$a = (x_1 \wedge x_2) + \dots + (x_{2n-1} \wedge x_{2n})$$

with  $x_1 \geq x_2 \geq \dots \geq x_{2n}$  in  $D$ .

Def. We call an embedding  $i: D \hookrightarrow B$  of distributive lattices tight if every  $a/b$  in  $B$  splits into  $a/b = x_1/y_1 + \dots + x_n/y_n$  with  $x_i, y_i \in D$ .

This  $i: D \hookrightarrow B(D)$  is tight.

Problem. Let  $X \hookrightarrow Y$  be a tight lattice embedding of MV algebras, and  $x > y, z > t$  in  $X$  with  $x \cdot y = z \cdot t$  in  $Y$ . Does this imply that  $x \cdot y = z \cdot t$  holds in  $X$ ?

Prop. 40. Let  $i: D \hookrightarrow B$  be a tight embedding of distributive lattices. If  $x \geq y \geq z$  with  $x, z \in D$ , there exist  $u, v \in D$  with  $u \geq v \geq z$  and  $y \vee u = x, y \wedge u = v$ .